

Solution: $\dim(W_1 + W_2) \leq \dim(V) = 10$ so 215

$$\dim(W_1) + \dim(W_2) - \dim(W_1 \cap W_2) \leq \dim(V) = 10 \quad \text{so}$$

$$7 + 6 - \dim(W_1 \cap W_2) \leq 10 \quad \text{so}$$

$$3 = 13 - 10 = 7 + 6 - 10 \leq \dim(W_1 \cap W_2).$$

Also, $W_1 \cap W_2 \subseteq W_1$ and $W_1 \cap W_2 \subseteq W_2$ so

$$\begin{aligned} \dim(W_1 \cap W_2) &\leq \text{Min}(\dim(W_1), \dim(W_2)) \\ &= \text{Min}(7, 6) = 6. \end{aligned}$$

Possibilities are: $3 \leq \dim(W_1 \cap W_2) \leq 6$

Unfinished business about det: 216

Recall: For $A = [a_{ij}] \in F_n^n$ we defined cofactors
 $A_{rs} = (-1)^{r+s} \det(M_{rs})$ for each $1 \leq r, s \leq n$.

We can get a new matrix $B = [A_{ij}] \in F_n^n$ from
these numbers, but we really want its transpose.

Def. The classical adjoint of $A \in F_n^n$ is

$$\text{adj}(A) = B^T = [A_{ji}].$$

Ex: $n=2$: $A = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$ so $B = [A_{ij}] = \begin{bmatrix} a_{22} & -a_{21} \\ -a_{12} & a_{11} \end{bmatrix}$

so $\text{adj}(A) = B^T = \begin{bmatrix} a_{22} & -a_{12} \\ -a_{21} & a_{11} \end{bmatrix}$. Note

$$A \text{adj}(A) = \begin{bmatrix} (a_{11}a_{22} - a_{12}a_{21}) & 0 \\ 0 & (a_{11}a_{22} - a_{12}a_{21}) \end{bmatrix} = (\det A) I_2$$

To see why $A \operatorname{adj}(A) = (\det A) I_n$ is true for (217) any $n \geq 2$, go back to cofactor expansions and prove this result:

Th. For $A = [a_{ij}] \in F^n$, with cofactors A_{rs} , we have

$$(a) \delta_{rs} \det(A) = \sum_{j=1}^n a_{rj} A_{sj} \quad \text{and}$$

$$(b) \delta_{rs} \det(A) = \sum_{i=1}^n a_{ir} A_{is} .$$

Pf. When $r=s$ this is the usual cofactor expansion theorem for computing $\det(A)$. When $r \neq s$ the right hand sides are cofactor expansions for the matrix obtained from A by replacing row (col.) s by row (col.) r , so those det's are 0 (two ident. rows or cols).

Cor. $\forall A \in F_n^n$, $A \operatorname{adj}(A) = \det(A) I_n = \operatorname{adj}(A) A$. (218)

Pf. The (r,s) -entry of $A \operatorname{adj}(A)$ is
 $\operatorname{Row}_r(A) \operatorname{Col}_s(\operatorname{adj}(A)) = \sum_{j=1}^n a_{rj} A_{sj} = \delta_{rs} \det(A)$.

The second equality, $\operatorname{adj}(A) A = (\det A) I_n$ comes from (b) of the Theorem above. \square

Th: (Cayley-Hamilton). For any $A \in F_n^n$, the char. poly. of A , $\operatorname{Char}_A(t) = \Delta(t) = \det(tI_n - A)$, is satisfied by A , that is, $\Delta(A) = O_n^n$.

Pf: (FAKE!) $\Delta(A) = \det(AI_n - A) = \det(O_n^n) = 0$.

WHY IS THIS NONSENSE?

If we write out the char. poly. of $A = [a_{ij}]$, (219)

$$\Delta_A(t) = \Delta(t) = \det(tI_n - A) = t^n + a_{n-1}t^{n-1} + \dots + a_1t + a_0$$

$$= \det \begin{bmatrix} t - a_{11} & \dots & -a_{1j} \\ \vdots & t - a_{22} & \vdots \\ -a_{ji} & \dots & t - a_{nn} \end{bmatrix} = \det [t\delta_{ij} - a_{ij}]$$

it clearly makes no sense to try to "plug in" $A \in F_n^n$ for t in the $n \times n$ matrix $tI_n - A$.

Let $B(t) = \text{adj}(tI_n - A)$, classical adjoint, whose entries are cofactors of $tI_n - A$, so are polys. in t of degree at most $n-1$. We can write

$$B(t) = B_{n-1}t^{n-1} + \dots + B_1t + B_0 \text{ for } B_i \in F_n^n \text{ constant}$$

matrices not involving t . Since $(tI_n - A)B(t) = \det(tI_n - A)I_n$ we have

$$(tI_n - A)(B_{n-1}t^{n-1} + \dots + B_1t + B_0) = (t^n + a_{n-1}t^{n-1} + \dots + a_1t + a_0)I_n$$

Multiply out the left side, group terms by 220 power of t and compare with corresponding terms on the right side:

$$\begin{aligned}
 & I_n B_{n-1} t^n + (I_n B_{n-2} - A B_{n-1}) t^{n-1} + (I_n B_{n-3} - A B_{n-2}) t^{n-2} \\
 & \dots + (I_n B_1 - A B_2) t^2 + (I_n B_0 - A B_1) t - A B_0 = \\
 & I_n t^n + a_{n-1} I_n t^{n-1} + a_{n-2} I_n t^{n-2} + \\
 & \dots + a_2 I_n t^2 + a_1 I_n t + a_0 I_n, \text{ gives:}
 \end{aligned}$$

$$\begin{aligned}
 B_{n-1} &= I_n \\
 B_{n-2} - A B_{n-1} &= a_{n-1} I_n \\
 B_{n-3} - A B_{n-2} &= a_{n-2} I_n \\
 &\vdots \\
 B_1 - A B_2 &= a_2 I_n \\
 B_0 - A B_1 &= a_1 I_n \\
 -A B_0 &= a_0 I_n
 \end{aligned}$$

so

$$\begin{aligned}
 A^n B_{n-1} &= A^n \\
 A^{n-1} B_{n-2} - A^n B_{n-1} &= a_{n-1} A^{n-1} \\
 A^{n-2} B_{n-3} - A^{n-1} B_{n-2} &= a_{n-2} A^{n-2} \\
 &\vdots \\
 A^2 B_1 - A^3 B_2 &= a_2 A^2 \\
 A B_0 - A^2 B_1 &= a_1 A \\
 -A B_0 &= a_0 I_n
 \end{aligned}$$

The sum of all terms on the left sides of [22]
These equations is O_n^n and the sum on the
right sides is $A^n + a_{n-1}A^{n-1} + \dots + a_1A + a_0I = \Delta(A)$,
so $\Delta(A) = O_n^n$ gives the result. \square

Cor. Let $L: V \rightarrow V$, $\dim(V) = n$, $A = {}_S[L]_S \in F_n^n$ for
some basis S of V , $\text{Char}_L(t) = \text{Char}_A(t) = \det(tI_n - A)$.
then $\text{Char}_L(L) = O_V^V \in \text{End}(V)$.

Discussion: In $\mathcal{U} = \text{End}(V) = \text{Lin}(V, V)$, the set
 $\{I_V, L, L^2, \dots, L^{n^2}\}$ is dep. since $\dim(\mathcal{U}) = n^2$ if
 $\dim(V) = n$. So there is some smallest $1 \leq m \in \mathbb{Z}$ s.t.
 $\{I_V, L, \dots, L^m\}$ is dep. but $\{I_V, L, \dots, L^{m-1}\}$ is indep.
Similarly; In F_n^n the set $\{I_n, A, \dots, A^{n^2}\}$ is dep.

since $\dim(F_n^n) = n^2$. Cayley-Hamilton Thm says $\{I_n, A, \dots, A^n\}$ is dep. so that smallest m s.t. $\{I_n, A, \dots, A^m\}$ is dep. but $\{I_n, A, \dots, A^{m-1}\}$ is indep. must be $m \leq n$.

Def. For $A \in F_n^n$, a minimal poly. for A is a poly. $m_A(t) \in F[t]$ which is monic (top coeff. 1) and is satisfied by A , and is the smallest possible degree. For $L: V \rightarrow V$ the corresponding minimal poly for L is denoted $m_L(t)$.

Recall discussion on page 70.

We have now answered some of those questions. For similar $B = P^{-1}AP$ what is the relation between

$m_A(t)$ and $m_B(t)$?

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Is $m_A(t)$ uniquely determined by A ?

Lemma: Let $B = P^{-1}AP$ and $f(t) \in F[t]$. Then

$$f(B) = P^{-1}f(A)P.$$

Pf. For each $0 \leq i \in \mathbb{Z}$, $B^i = P^{-1}A^iP$ can be proven by induction on i . For $i=0$, $B^0 = I_n = P^{-1}I_nP = P^{-1}A^0P$.

Assuming the equation is true for some $0 \leq i \in \mathbb{Z}$,

$$B^{i+1} = (P^{-1}AP)^{i+1} = (P^{-1}AP)^i(P^{-1}AP) = (P^{-1}A^iP)(P^{-1}AP)$$

$$= P^{-1}A^i(P P^{-1})AP = P^{-1}A^i I_n AP = P^{-1}A^i A P = P^{-1}A^{i+1}P.$$

$$\text{So } f(t) = \sum_{i=0}^m c_i t^i \text{ so } f(B) = \sum_{i=0}^m c_i B^i = \sum_{i=0}^m c_i (P^{-1}A^iP)$$

$$= P^{-1} \left(\sum_{i=0}^m c_i A^i \right) P = P^{-1}f(A)P. \quad \square$$

So if $f(A) = 0_n^n$ then $f(B) = P^{-1}f(A)P = 0_n^n$, 224
says similar matrices are satisfied by the same
polys.

Th: Let $A \in F_n^n$ and let $\mathcal{I}_A = \{f(t) \in F[t] \mid f(A) = 0_n^n\}$.
Then $\mathcal{I}_A \trianglelefteq F[t]$ is an ideal in the ring $F[t]$ and
 $\mathcal{I}_A = \{g(t)m_A(t) \in F[t] \mid g(t) \in F[t]\} = (m_A(t))$ is
the (principal) ideal of all multiples of $m_A(t)$,
the unique monic poly of minimal degree in \mathcal{I}_A .
Pf. It is clear that \mathcal{I}_A is closed under $+$ and \cdot , so
it is a subspace of $F[t]$, and \mathcal{I}_A absorbs products,
 $\forall f(t) \in \mathcal{I}_A, \forall g(t) \in F[t], f(t)g(t) \in \mathcal{I}_A$, so it is an
ideal. The Euclidean Algorithm in $F[t]$ (see page 66)

implies that any ideal of $F[t]$ is principal, (225)
but here we apply it only to \mathcal{I}_A . Let $m_A(t)$ be a
poly. of minimal degree in \mathcal{I}_A . Multiplying by a scalar
we can assume it is monic. Then $\forall f(t) \in \mathcal{I}_A$,

$\exists q(t), r(t) \in F[t]$ s.t. $f(t) = q(t)m_A(t) + r(t)$
where either $r(t) = 0$ or $\deg(r(t)) < \deg(m_A(t))$.

Evaluation of both sides at $t = A$ gives

$0_n^n = f(A) = q(A)m_A(A) + r(A)$ and $m_A(A) = 0_n^n$
so $0_n^n = r(A)$. If $\deg(r(t)) < \deg(m_A(t))$ we get
a contradiction to $m_A(t)$ being the smallest degree
poly. in \mathcal{I}_A so $r(t)$ must be the zero poly. The
uniqueness of $m_A(t)$ follows because any two polys
of min. degree in \mathcal{I}_A must divide each other.
They can only differ by a scalar multiple. Monic \Rightarrow Unique. \square