

Thus, $x_s - s$ is a difference of two commuting (95)
s.s. endo's of V , so it is s.s., and
 $n - x_n$ is a diff. of two comm. nilp. endo's
of V , so it is nilp.

The only operator in $\text{End}(V)$ which is both s.s.
and nilp. is the zero operator, 0_V .
So $s = x_s$ and $n = x_n$ gives uniqueness. \square

Applications: For $\dim(V) < \infty$ let $L = \mathfrak{gl}(V)$
and its rep'n $\text{ad}: L \rightarrow \text{End}(L)$.
We have seen that $x \in \mathfrak{gl}(V)$ nilp. $\Rightarrow \text{ad}_x$ nilp.

Lemma: If $x \in \mathfrak{gl}(V)$ is s.s. then ad_x is s.s.

Pf. Let $S = \{v_1, \dots, v_n\}$ be a basis of V such that

$x \cdot v_i = a_i \cdot v_i$ so the matrix rep'ing x w.r.t. S is $\text{diag}(a_1, \dots, a_n)$. The "std. basis" of $\mathfrak{gl}(V)$ associated with S is the set

$\{E_{ij} \in \text{End}(V) \mid 1 \leq i, j \leq n\}$ where $E_{ij}(v_k) = \delta_{jk} v_i$ for $1 \leq k \leq n$. The $n \times n$ matrix representing

E_{ij} w.r.t. S is $E_{ij} \in F^n$ which has 1 in its (i, j) th entry, 0's everywhere else.

To compute $\text{ad}_x(E_{ij})$ we apply it to any v_k :

$$(x \circ E_{ij} - E_{ij} \circ x) v_k = \delta_{jk} a_i v_i - a_k \delta_{jk} v_i$$

$$= (\delta_{jk} a_i - \delta_{jk} a_j) v_i = (a_i - a_j) \delta_{jk} v_i = (a_i - a_j) E_{ij} v_k$$

so $\text{ad}_x(E_{ij}) = (a_i - a_j) E_{ij}$ shows ad_x is s.s. \square

Lemma A: Let $X \in \text{End}(V)$ for $\dim(V) < \infty$, [97] and let $X = X_s + X_n$ be its Jordan decomposition. Then $\text{ad}_X = \text{ad}_{X_s} + \text{ad}_{X_n}$ is the J-decomp. of $\text{ad}_X \in \text{End}(\text{End}(V))$.

P.f. We know ad_{X_s} is s.s. and ad_{X_n} is nilp. and $[\text{ad}_{X_s}, \text{ad}_{X_n}] = \text{ad}[X_s, X_n] = \text{ad}_0 = 0$.

The uniqueness of J-decomp. from Prop. (a) gives the result. \square

Lemma B. Let $(\mathcal{A}, *)$ be any fin. dim'l F-alg.

Then $\text{Der}(\mathcal{A}) = \left\{ \delta \in \text{End}(\mathcal{A}) \mid \begin{array}{l} \delta(x*y) = \delta(x)*y + \\ x*\delta(y), \forall x, y \in \mathcal{A} \end{array} \right\}$
 and $\forall \delta \in \text{Der}(\mathcal{A})$, the s.s. and nilp. parts of δ , δ_s and δ_n , are each in $\text{Der}(\mathcal{A})$.

Pf. Let $\delta \in \text{Der}(A)$, $\sigma = \delta_s$, $\nu = \delta_n \in \text{End}(A)$ 198
 so $\delta = \sigma + \nu$. We will show $\sigma \in \text{Der}(A)$, and
 that will imply $\nu = \delta - \sigma \in \text{Der}(A)$.

For $a \in F$ let $\mathcal{Q}_a = \{x \in A \mid (\delta - aI)^k x = 0 \text{ for}$
 some k (dep'ing on x) $\}$.

Then $A = \bigoplus \mathcal{Q}_a$
 where the sum is over those $a \in F$ s.t. a is an
 e. value for δ (or for σ) and $\sigma|_{\mathcal{Q}_a} \rightarrow \mathcal{Q}_a$ so
 that the restriction $\sigma|_{\mathcal{Q}_a} = aI_{\mathcal{Q}_a}$ is scalar mult
 by $a \in F$.

For any $a, b \in F$, $\forall x, y \in A$ we have

$$(*) (\delta - (a+b)I)^n (x * y) = \sum_{i=0}^n \binom{n}{i} (\delta - aI)^{n-i} x * (\delta - bI)^i y$$

for $0 \leq n \in \mathbb{Z}$, by induction on n .

The $n=0$ case is trivial: $x * y = x * y$. [99]

$$\text{For } n=1: (\delta - (a+b)I)(x * y) = \delta(x * y) - (a+b)(x * y)$$

$$= \delta(x) * y + x * \delta(y) - a(x * y) - b(x * y)$$

$$= ((\delta - aI)(x)) * y + x * ((\delta - bI)y).$$

Inductive Step: Exercise.

(Did you use assoc. of $*$ in \mathcal{A} at any point?)

Then (*) implies $\mathcal{A}_a * \mathcal{A}_b \subseteq \mathcal{A}_{a+b}$. (Check!)

If $x \in \mathcal{A}_a$, $y \in \mathcal{A}_b$ then $\sigma(x * y) = (a+b)(x * y)$
since $x * y \in \mathcal{A}_{a+b}$ where σ acts as scalar $a+b$.

But then $\sigma(x * y) = \sigma(x) * y + x * \sigma(y) = (ax) * y + x * (by)$
and $\mathcal{A} = \bigoplus \mathcal{A}_a$ means $\sigma \in \text{Der}(\mathcal{A})$. \square

Cartan's Criterion (for solvability of L) (100)

Note: If $[L, L]$ is nilp. then L is solvable
since $L^{(0)} = [L, L] = L'$ and $L^{(i)} \subseteq L^i$ for $i \geq 0$.

Also: $[L, L]$ is nilp. iff $\forall x \in [L, L]$,
 $ad_x: [L, L] \rightarrow [L, L]$ is nilp. by Engel's Th.

So we first develop a trace criterion for
 $x \in \text{End}(V)$ to be nilp.

Lemma: Let $A \subset B \subset \mathfrak{gl}(V)$ be subspaces
with $\dim(V) < \infty$. Let $M = \{x \in \mathfrak{gl}(V) \mid [x, B] \subset A\}$.

Suppose $x \in M$ satisfies $\text{Tr}(xy) = 0, \forall y \in M$.
Then x is nilp.

P.F. Let $X = S + N$ be the Jordan decomposition [10] of X , so $S = X_S$, $N = X_N$ in previous notation. Let $S = \{v_1, \dots, v_m\}$ be a basis of V s.t. matrix of S is $\text{diag}(a_1, \dots, a_m)$, that is, $S \cdot v_i = a_i v_i$. Let $E = Qa_1 + \dots + Qa_m = Q\text{-span of } \{a_1, \dots, a_m\} \subseteq F$ considered as a vector space over $Q \subseteq F$ since $\text{char}(F) = 0$. To show X is nilp. we want to show $S = 0$, that is, $E = 0$. $\dim_Q(E) < \infty$ so $\dim(E) = \dim(E^*)$, dual space, so we want to show that any lin. map $f: E \rightarrow Q$ is zero. $\forall f \in E^*$, let $\gamma \in \text{gl}(V)$ be defined by $\gamma \cdot v_i = f(a_i) v_i$ for $1 \leq i \leq m$, so the matrix rep'ing γ w.r.t. S is $\text{diag}(f(a_1), \dots, f(a_m))$. As before, let $\{E_{ij} \in \text{gl}(V) \mid 1 \leq i, j \leq m\}$ be the

basis of $\mathfrak{gl}(V)$ s.t. $E_{ij}(v_k) = \delta_{jk} v_i$ and so 102

$\text{ads}(E_{ij}) = (a_i - a_j) E_{ij}$ and $\text{ady}(E_{ij}) = (f(a_i) - f(a_j)) E_{ij}$
for $1 \leq i, j \leq m$ since s and y are semisimple.

Let $r(T) \in F[T]$ be a poly with $r(0) = 0$ s.t.

$r(a_i - a_j) = f(a_i) - f(a_j)$ for all $1 \leq i, j \leq m$.

Such a poly. is found by Lagrange interpolation.

If $a_i - a_j = a_k - a_l$ then lin. of f gives the
consistent value $f(a_i - a_j) = f(a_k - a_l)$. Then,
evaluation of poly. $r(T)$ at $T = \text{ads}$ yields the

endomorphism $r(\text{ads}) \in \text{End}(\mathfrak{gl}(V))$ such that

$$r(\text{ads})(E_{ij}) = r(a_i - a_j) E_{ij} = (f(a_i) - f(a_j)) E_{ij}$$

$$= \text{ady}(E_{ij}) \text{ so } r(\text{ads}) = \text{ady}.$$