## Lecture 3 (Conclusion of Theorem 3) 8 November 2019 Lecture notes by Michael Gottstein

We finish the proof of Theorem 3. First, a few new definitions.

**Definition 1.** A *unicyclic graph* is a connected graph with exactly one circle. A 1-*tree* is a tree with or without a half edge. A *pseudotree* is a 1-tree or a unicyclic graph. A *pseudoforest* is a graph whose components are pseudotrees.

**Definition 2.** A gain graph is *contrabalanced* if every circle is unbalanced; i.e., it has no balanced circles.

**Theorem 3.** Let  $\Phi$  be a  $K^{\times}$ -gain graph, Let M be the matroid on  $E(\Phi)$  that corresponds to  $\mathscr{M}(\mathscr{A}([\Phi]))$ . The independent sets in M are (the edge sets of) the contrabalanced pseudoforests in  $\Phi$ .

Recall that Lemma 3 says an edge set each of whose components is a tree, an unbalanced unicycle, or a tree with a single half edge is independent in the matroid M of the gain graph  $\Phi$  that corresponds to the hyperplane arrangement  $\mathscr{A}[\Phi]$  over a field K. I.e., any contrabalanced pseudoforest is independent in M.

So we must prove every independent edge set is a contrabalanced pseudoforest. We'll prove the contrapositive:

**Lemma 4.** The dependent sets are the edge sets that have a single component with at least 2 circles, at least 2 half edges, or at least one circle and one half edge.

*Proof.* Since a set containing a dependent set is dependent we only need to prove that a connected edge set that contains, two circles, two half edges, or a circle and half edge is dependent in M.

Suppose S is such an edge set. If S contains a balanced circle then we know it is dependent by Case 1 of the previous treatment of a circle. So we may assume every circle in S is unbalanced.

Case 1: S contains an unbalanced circle or half edge  $C_1$ , and one other one  $C_2$ , that share at most one vertex. There must be a minimal path P connecting  $C_1$  and  $C_2$  (see figure 2)

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By Case 2 of circles or by definition of a half edge hyperplane h(e),  $C_1$  forces  $x_i = 0$  and  $C_2$  forces  $x_i = 0$ .

Indeed, write  $P = w_0 e_1 w_1 \dots w_{l-1} e_l w_l$  so that in  $\mathscr{A}[\Phi]$ ,

$$x_{w_l} = x_{w_{l-1}}\varphi_{l-1,l}(e_l) = \dots = x_{w_0}\varphi(P).$$

But  $x_{w_0} = 0$  from  $C_1$ , therefore  $x_{w_l} = 0$ .

We will show that some  $e \in S$  has the property that

$$h(e) \supseteq \bigcap_{f \in S \setminus e} h(f).$$

(Every  $e \in S$  has that property, but we don't need that.)

Pick e to be an edge in  $C_2$  incident with  $v_j$  so  $e = e_{jk}$  or the half edge  $e_j$ . With the equations  $x_k = x_j \varphi_{jk}(e)$  for a circle and  $x_j = 0$  a half edge.

Then we consider the equations separately.

In the half edge case  $x_j = 0$  (consider  $C_1$  and P) in  $\bigcap h(C_1 \cup P) = \bigcap h(S \setminus e_j)$ , therefore h(S) is a dependent set of hyper planes.

In the circle case  $x_j = 0$  and  $x_k = 0$  (consider  $C_1$  and the path  $P \cup (C_2 \setminus e)$ ), then  $x_k = x_j \varphi_{jk}(e)$  is satisfied (with  $x_k = x_j = 0$ ) by  $\bigcap h(S \setminus e)$ ; therefore h(S) is a dependent set of hyperplanes.

This solves the problem when S contains two unbalanced circles/half edges with at most one vertex in common.

Case 2: Now we will consider the case where S contains two circles with at least two common vertices.

We treat first the case where S is a theta graph. We have 3 unbalanced circles, say  $C_{12} = P_1 P_2^{-1}$ ,  $C_{23} = P_1 P_3^{-1}$ ,  $C_{13} = P_1 P_3^{-1}$  given by the internally disjoint paths of the theta graph,  $P_1, P_2, P_3$ , which all start and end at  $v_i$  and  $v_j$ , respectively. Then  $C_{12}$ , being unbalanced, implies  $x_i = x_j = 0$ . Let  $e = e_{jk}$  be the edge in  $p_3$  at  $v_j$  Then  $C_{12} \cup (P_3 \setminus e)$  implies  $x_k = 0$  thus  $x_k = 0$  in  $\bigcap h(C_{12} \cup (P_3 \setminus e)) =$  $\bigcap h(S \setminus e)$ . Similarly,  $x_j = 0$ . Therefore  $h(e) \supseteq h(S \setminus e)$ , so the theta graph is dependent in M. So if S contains a theta graph, it is dependent.

Note that we have been assuming S is connected. Let  $\xi(S)$  be the cyclomatic number of S, defined as the smallest number of edges that when deleted leave a tree spanning V(S), i.e., the minimum number of edges whose deletion leave a forest (therefore a tree). We have  $\xi(S) > 1$  because  $C_1$  and  $C_2$  exist. If  $\xi(S) = 2$  then S contains the theta graph  $C_1 \cup C_2$  (because those two circles have at least two common vertices), so we are done. If  $\xi(S) > 2$  then we can delete  $\xi(S) - 2$  edges to get a connected subgraph S' with  $\xi(S') = 2$ , which contains a theta graph or handcuff (by easy graph theory) and is therefore dependent.

So we have finished the proof modulo some graph-theoretic detail.

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