

GAIN GRAPHS AND HYPERPLANE ARRANGEMENTS  
LECTURE 4 (THE FRAME MATROID)  
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In the last lecture, we analyzed the matroid  $\mathcal{M}(\mathcal{A}[\Phi])$  in terms of the independent sets of  $\Phi$ . Theorem 1 showed that the independent sets are contrabalanced pseudoforests.

**Definition 1.** Let  $\mathbf{F}(\Phi)$  be this matroid on a  $K^\times$ -gain graph  $\Phi$ , with independent sets given by Theorem 1. This is the *frame matroid* of  $\Phi$ .

**Theorem 1.** For the gain graph  $\Phi$ , let  $S \subseteq E(\Phi)$ . Then,

- (1)  $rk(S) = n - b(S)$ .
- (2) The circuits of  $F(\Phi)$  belong to one of the following three categories:
  - (a) *Balanced circles and loose edges.*
  - (b) *Contrabalanced handcuffs.*
  - (c) *Contrabalanced theta graphs.*
- (3) The closure of  $S$  is given by

$$cl(S) = (E:V_0(S)) \cup bcl(S:V_0(S)^c).$$

- (4)  $S$  is closed if  $S$  is equal to its closure (i.e. the union of unbalanced components is an induced subgraph of  $\Phi$  and each balanced component is balance-closed)

To explain the notation we give some additional definitions.

**Definition 2.** For  $S \subseteq E$ , we define  $b(S)$  to be the number of balanced components of the spanning subgraph  $(V, S)$ . Isolated vertices are included as balanced components because they are balanced. However, loose edges are not included as they are not considered components.

We define  $V_0(S)$  to be the set of vertices that are contained in unbalanced components of  $S$ .

**Definition 3.** For  $X \subseteq V$  and  $S \subseteq E$ , the *induced subset of  $S$  on  $X$* , denoted by  $S:X$ , is defined as the set  $\{e \in S : \text{all endpoints of } e \text{ are in } X\}$ . Loose edges are included in the induced subset, as they do not contain a vertex that is not in  $X$ .

**Definition 4.** The *balance-closure of  $S$* , denoted by  $bcl(S)$ , is given by

$$S \cup \{e \in E : \exists \text{ balanced circle } C \text{ such that } e \in C \subseteq S \cup \{e\}\} \\ \cup \{\text{loose edges}\},$$

which is equal to

$$\begin{aligned} S \cup \{e \in E : \exists \text{ path } P \text{ in } S \text{ joining the endpoints of } e \\ \text{such that } P \cup \{e\} \text{ is a balanced circle}\} \\ \cup \{\text{balanced loops}\} \cup \{\text{loose edges}\}. \end{aligned}$$

In both sets the loose edges could be omitted if we consider loose edges to be balanced circles.

An edge set  $S$  is *balance-closed* if  $S = \text{bcl}(S)$ .

Before we begin the proof of the theorem, there are two notes about the balance-closure of  $S$ . First, we do not yet know if the balance-closure satisfies all three conditions to be an abstract closure. Clearly,  $S \subseteq \text{bcl}(S)$  and balance-closure is weakly increasing (i.e., adding edges to  $S$  cannot remove an edge from  $\text{bal}(S)$ ). However, it may not be true that  $\text{bcl}(\text{bcl}(S)) = \text{bcl}(S)$ . Secondly, the balance-closure of a balanced set must be balanced, because it contains no unbalanced circles or half-edges. This requires proof.

**Lemma 1.** *If  $S$  is balanced, then  $\text{bcl}(S)$  is balanced.*

*Proof.* Switch so  $S$  has all identity gain. Then an edge added to  $S$  by balance-closure must also have identity gain. It follows that  $\text{bcl}(S)$  is balanced.  $\square$

We begin the proof of the theorem with part (2).

*Proof of Part (2).* Let  $D$  belong to one of the categories described in this part and let  $e \in D$ . From Theorem 1, we know that  $D$  is dependent and that  $D \setminus e$  is independent. Independence comes from the fact that if we were to remove any edge from  $D$ , this would result in a contrabridged pseudoforest. Therefore  $D$  is a circuit.

Conversely, let  $D$  be a circuit. Then  $D$  is dependent. By Theorem 1,  $D$  must contain a balanced circle or two unbalanced circles/half-edges that are in the same component of  $D$  (since if they were in different components, they would not make  $D$  dependent). In Theorem 1 we already showed that this subset  $D'$  of  $D$  is dependent and is either a balanced circle or a contrabridged handcuff or theta graph. By minimality,  $D = D'$ .  $\square$

*Proof of Part (1).* From the definition of  $\mathbf{F}(\Phi)$  (as corresponding to  $\mathcal{M}(\mathcal{A}[\phi])$ ) we have  $\text{rk}(S) = \text{rk}(\bigcap h(S))$ , as the right-hand side is the rank function in the hyperplane arrangement. This is then equal to  $\text{codim} \bigcap h(S)$ . We know that  $\text{rk}(S) = \text{rk}(I)$  for any maximal contrabridged pseudoforest  $I$  because it is a maximal independent set. Thus we also have  $\text{rk}(S) = \text{rk}(I) = \text{codim} \bigcap h(I)$ .

The following is a diagram of the structure of  $I$ .

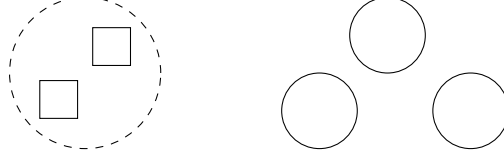


FIGURE 1. Structure of  $I$ . Each box represents an unbalanced component  $R_i$  of  $I$  and each circle represents a balanced component  $I_i$  of  $I$ . The dashed circle represents  $V_0(S)$ .

We can break up  $I$  into its unbalanced components and its balanced ones. Collectively, the unbalanced components are  $I:V_0(S)$ . For the balanced components, each component must be a tree because  $I$  is independent. (We are able to ignore loose edges as they correspond to the degenerate hyperplane.)

For an unbalanced component  $R_i$ , we found in the study of circles and the proof of Theorem 1 that  $h(R_i) \implies x_j = 0$  for all  $v_j \in V(R_i)$ . Hence, we have a contrabalanced 1-tree. Therefore  $\text{codim} \bigcap h(R_i)$  will be equal to the number of coordinates of the hyperplane equations in  $h(R_i)$ . However, this is just the number of vertices of  $R_i$ . Therefore  $\text{rk}(R_i) = |V(R_i)|$ .

Next, we consider a balanced component  $I_i$ . It is a tree, by Theorem 1. All coordinates of  $\bigcap h(I_i)$  are determined through the equations of the tree by fixing any one coordinate arbitrarily (see Figure 2). Therefore  $\text{codim} \bigcap h(I_i) \geq |V(I_i)| - 1 = |E(I_i)|$ .

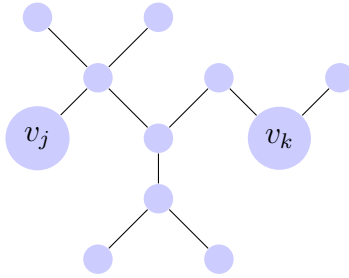


FIGURE 2. A possible structure for  $I_i$ . By fixing  $x_j$ , we can find  $x_k$  in terms of  $x_j$ .

Since  $\text{codim} \bigcap h(I_i) \leq |E(I_i)|$  because the codimension cannot be greater than the number of hyperplanes of  $h(I_i)$ , therefore  $\text{codim} h(I_i) = |V(I_i)| - 1$ .

Now, since each  $h(R_i)$  and each  $h(I_i)$  uses a different set of coordinates, we can write  $\text{rk}(I)$  as  $\sum_i \text{rk}(R_i) + \sum_i \text{rk}(I_i)$  (see the previous lemma about direct sums). Using the formulas we have just constructed, this becomes  $\sum_i |V(R_i)| + \sum_i (|V(I_i)| - 1) = |V(I)| - b(I)$ . Each vertex is included in either some  $R_i$  or some  $I_i$ , so this is  $n - b(I)$ .

For the proof of Part (1) to be complete, it remains to be shown that the vertex sets of the balanced components of  $I$  are those of the balanced components of  $S$ ; in other words, that  $I$  restricted to a balanced component of  $S$  is connected.  $\square$

**Lemma 2.** *Suppose  $I$  is a maximal independent set in  $S \subseteq E$ . Let  $S$  have unbalanced components  $U_1, U_2, \dots$  and balanced components  $B_1, B_2, \dots$ . Then  $I \cap B_i$  is a tree spanning  $V(B_i)$  and  $I \cap U_i$  is a disjoint union of contrabalanced 1-trees.*

*Proof.* To appear next time.  $\square$