## Gain Graphs and Hyperplane Arrangements LECTURE 5 13 November 2019 Lecture Notes by Nicholas Lacasse

In the last lecture we ended with the statement of the following lemma which is necessary to complete the proof of Lemma 1 of Lecture 4. We begin with proving this lemma.

**Lemma 1.** Let  $\Phi$  be a  $K^{\times}$ -gain graph,  $S \subseteq E(\Phi)$ , and I a maximal independent set contained in S. Then  $b(I) = b(S)$ .

*Proof.* Suppose S has balanced components  $B_1, B_2, \ldots$  and unbalanced components  $U_1, U_2, \ldots$  To prove that I and S have the same number of balanced components we will show two things:

- (1)  $I \cap E(B_i)$  is a spanning tree of  $B_i$  (and therefore balanced),
- (2)  $I \cap E(U_i)$  is a contrabalanced spanning 1-forest (and therefore has no balanced components).

For (1), notice that  $I \cap B_j$  is a forest. This is because  $B_j$  is balanced,  $I \subseteq B$ , and I is independent. (If  $I \cap B_j$  were not a forest there would be a balanced circle contained in I, but we know balanced circles are dependent. So  $I \cap B_j$  is a forest.) Suppose it is disconnected. Then we can use an edge of  $B_j$  to join two trees of  $I \cap B_j$  into one tree, since  $B_j$  is connected. Therefore,  $I \cap B_j$  was not maximal, and thus I was not maximal. But we chose  $I$  to be maximal, so this is a contradiction. Thus  $I \cap B_j$  is a tree, which spans  $B_j$  because an isolated vertex is a (very small) tree, so it falls under the previous argument.

For (2), notice that the only possible balanced components of  $I \cap U_i$ are trees. By way of contradiction, suppose  $I \cap U_j$  has a tree component, T. If T is not the only component of  $I \cap U_j$  then we can add an edge of  $U_j$  to I in order to connect T to another component U of  $I \cap U_j$  (this other component of  $I \cap U_j$  will necessarily be either a tree, a tree with a single half-edge, or an unbalanced unicycle because I is a contrabalanced pseudoforest), giving a larger independent set (because connecting  $T$  to a component which is either a tree, a tree with a single half-edge, or an unbalanced unicycleU will just yield a larger component of the same type as  $U$ , preserving the contrabalanced pseudoforest). Since we chose  $I$  to be maximally independent, this is a contradiction. So T must be the only component of  $I \cap U_j$ . Then either  $U_j$  has a half-edge  $e$ , in which case  $T \cup e$  is a larger independent set, or  $U_i$  has an unbalanced circle C, which can be extended to a unicycle U that spans  $U_j$  (since  $U_j$  is connected). We can replace T by  $U$  to create an independent set larger than  $I$ . This is a contradiction because maximal independent sets all have the same size (by matroid theory). So  $I \cap U_j$  has no tree components. Therefore  $I \cap U_j$  has no balanced components. So in  $I$ , we get exactly one balanced component for each  $B_j$  and no others. Therefore,  $b(I) = b(S)$ .

N.B. Once we have established the independent sets of the frame matroid  $\mathbf{F}(\Phi)$  using  $\mathscr{M}(\mathscr{A}[\Phi])$ , we can define  $\text{rk}(S) = \max\{\#I \mid I \subseteq I\}$ S, I independent for  $S \subseteq E(\Phi)$  entirely in terms of the gain graph. We never need to refer back to  $\mathscr{A}[\Phi]$ . This permits a vast generalization. Take any biased graph  $(\Gamma, \mathscr{B})$ , possibly from a gain graph and possibly not, and define  $S \subseteq E$  to be independent if it is a contrabalanced pseudoforest.

**Theorem 2.** Define  $\mathbf{F}(\Gamma, \mathcal{B}) := ((\Gamma, \mathcal{B}), \mathcal{I})$  where  $\mathcal{I}$  is the set of all  $S \subseteq E(\Gamma)$  which induce contrabalanced pseudoforests. Then  $\mathbf{F}(\Gamma, \mathcal{B})$  is a matroid on E.

*Proof.* This proof will be omitted or postponed.

Given Theorem 2, our proof that  $rk(S) = n - b(S)$  for  $S \subseteq E$  and our proof of the circuits of  $\mathbf{F}(\Phi)$  both apply to  $\mathbf{F}(\Gamma,\mathscr{B})$ .

The next step in the proof of Theorem ?? is to establish the closure operator.

**Lemma 3.**  $cl_F(S) = (E: V_0(S)) \cup bel(S: V_0(S)^c)$ .

For this we will establish a helpful lemma.

**Lemma 4.** Let  $\Phi$  be a gain graph and let  $R \subseteq E$  be balanced and connected. Then  $\text{bel}(R)$  (aside from the balanced loops and loose edges) is the maximal balanced subset of  $E$  that contains  $R$  and has vertex set contained in  $V(R)$ .

*Proof.* Since R is balanced we can switch  $\Phi$  so it has all identity gain. Then  $\text{bel}(R) = \{e \in E: V(R) \mid \varphi(e) = \epsilon\}$  together with all balanced loops and loose edges.

For an arbitrary edge set  $R$  we have a reduction to components, which is of most use when  $R$  is balanced.

**Lemma 5.** Let  $\Phi$  be a gain graph and  $R \subseteq E$ . Let  $\Omega$  be the set of components of R. Then  $\mathrm{bel}(R) = \bigcup_{C \in \Omega} \mathrm{bel}(C)$ .

Proof. An exercise. □

Balance-closure differs from closure in that bcl never joins unbalanced components of R, but cl can join them.